

Automatic Program Reoptimization Support in LLVM ORC JIT

by Sunho Kim







Undergrad student from UC San Diego

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 - $\circ~$ Which is what this talk will be about

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- Compile with -O2 for only "hot" functions
 - $\circ\,$ The compilation time of -O0 or -O1 is faster than -O2 in general
- Runtime profile guided optimization
 - $\circ\,$ De-virtualization, instruction reordering, and other types of PGOs in ORC JIT
- Scientific computing (CERN)
 - Use high precision floating point for early iterations and use low precision floating point in later iterations for places that matter

REVIVING FEATURE FROM 2003?

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Ilvm-project / Ilvm / docs / HistoricalNotes / 2003-06-25-Reoptimizer1.txt

🝘 ddunbar [typo] An LLVM. 🚥

Blame 137 lines (105 loc) · 5.88 KB Code Wed Jun 25 15:13:51 CDT 2003 1 First-level instrumentation 3 -----We use opt to do Bytecode-to-bytecode instrumentation. Look at 6 back-edges and insert llvm_first_trigger() function call which takes no arguments and no return value. This instrumentation is designed to be easy to remove, for instance by writing a NOP over the function 9 10 call instruction. 11 12 Keep count of every call to llvm_first_trigger(), and maintain 13 counters in a map indexed by return address. If the trigger count exceeds a threshold, we identify a hot loop and perform second-level 14 15 instrumentation on the hot loop region (the instructions between the 16 target of the back-edge and the branch that causes the back-edge). We 17 do not move code across basic-block boundaries. 18 19 20 Second-level instrumentation 21 -----22 23 We remove the first-level instrumentation by overwriting the CALL to 24 llvm_first_trigger() with a NOP. 25

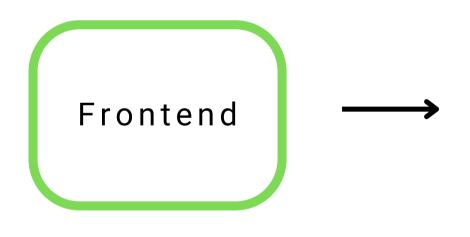
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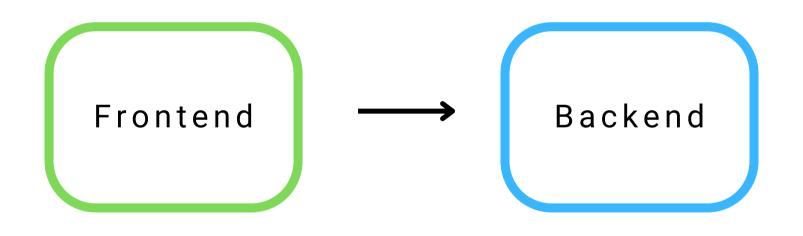
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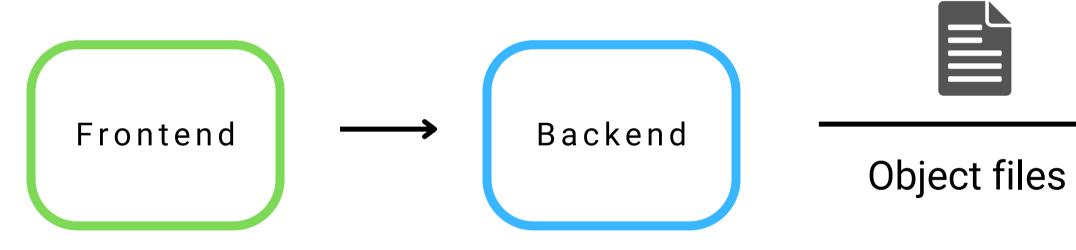
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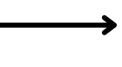
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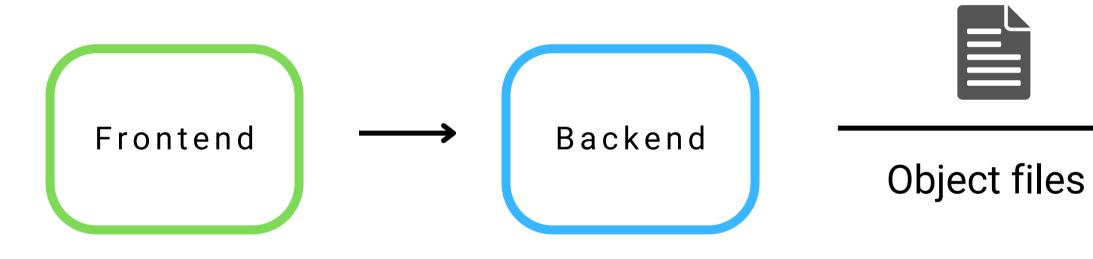
Quite different but has the same name :)





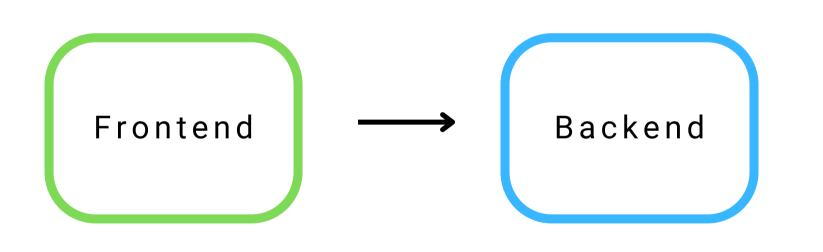




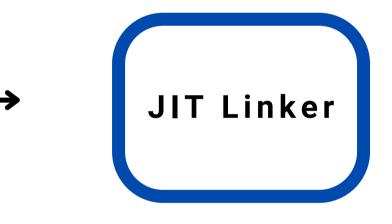




JIT execution pipeline in LLVM



Object files (in memory)



JIT execution pipeline in LLVM



- Share a huge portion of pipeline with AOT
- Fewer breakage by LLVM internal code changes



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 - Frontend AST or IR module will start compiling when a function defined by it is called in runtime.

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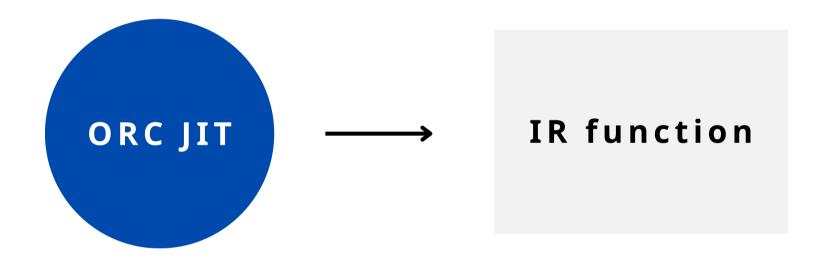
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- Multi-thread, remote process, speculative compilation ...

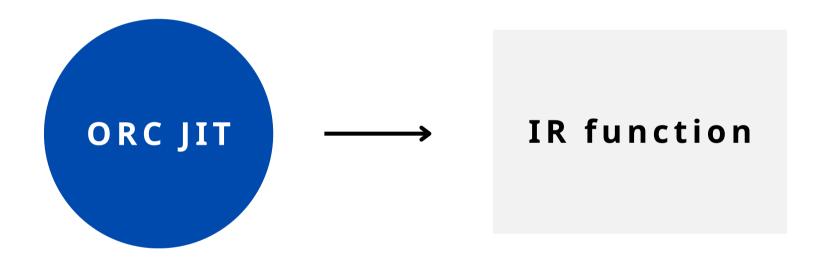
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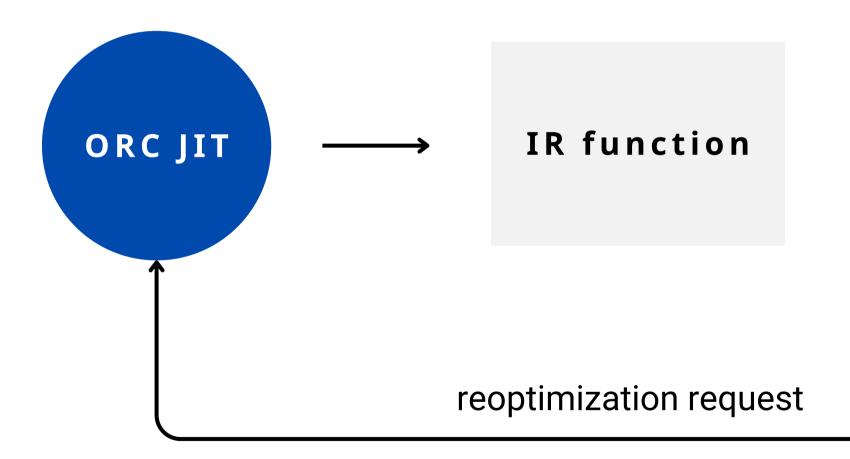


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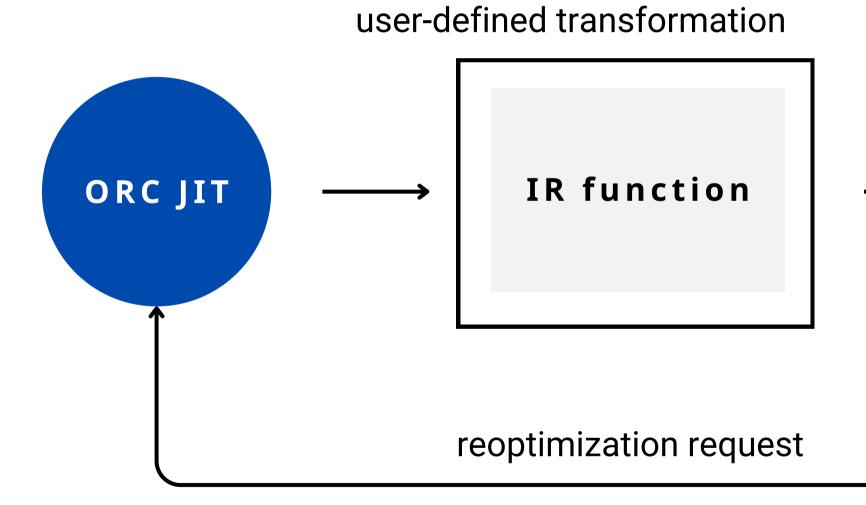
Binary code

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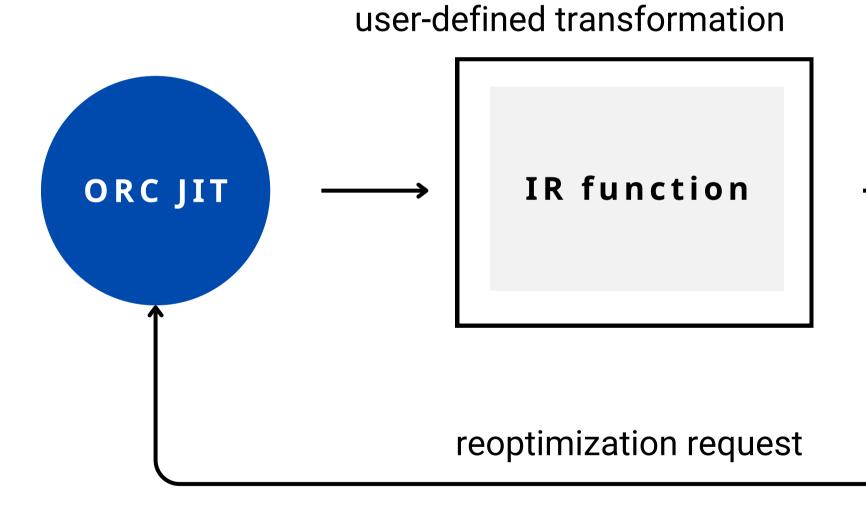
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Reoptimized Binary Code

• LLLayerJIT

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std::unique_ptr<LLLayerJIT> Jit;

Jit->addLayer(ReOptLayer);

Jit->addLayer(std::make_unique<LLIRPartitionLayer>());

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• LLLayerJIT

std::unique_ptr<LLLayerJIT> Jit;

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it IR module Add lazy-compilation layer

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```
static Error reoptimizeBasic(ReOptimizeLayer &Parent, ReOptMaterializationUnitID MUID,
   unsigned CurVerison, ResourceTrackerSP OldRT, ThreadSafeModule &TSM) {
 TSM.withModuleDo([&](llvm::Module &M) {
    // Do some re-optimization based on profile data
 });
 return Error::success();
auto ReOptLayer = std::make_unique<LLReOptimizeLayer>(ES, RSManager);
ReOptLayer->setReOptimizeFunc(reoptimizeBasic);
```

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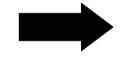
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- Called to add instrumentation code to the "first version" of the functions.
- Default is "reoptimizelfCallFrequent" which requests re-optimization when call count is high.

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 entry:
 ret i32 5
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define i32 @hi() { entry: ret 132 5 }



define di32 @hi() { entry: %cnt_is_20 = icmp eq i64 %0, 20 %cnt_plus_one = add i64 %0, 1 reoptimize_request: br label %4 return: ret i32 5 3

```
%cnt = load i64, ptr @__orc_reopt_counter, align 8
store i64 %cnt_plus_one, ptr @__orc_reopt_counter, align 8
br i1 %cnt_is_20, label %reoptimize_request, label %return
```

```
call void @__orc_rt_reoptimize(i64 3, i32 0)
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call void @__orc_rt_reoptimize(i64 3, i32 0)

Example: do -O2 optimization if function was called more than 10

```
static Error reoptimizeToO2(ReOptimizeLayer &Parent, ReOptMaterializationUnitID MUID,
   unsigned CurVerison, ResourceTrackerSP OldRT, ThreadSafeModule &TSM) {
 TSM.withModuleDo([&](llvm::Module &M) {
   auto PassManager = buildPassManager();
   PassManager.run(M);
 });
 return Error::success();
ReOptLayer->setReOptimizeFunc(reoptimizeToO2);
ReOptLayer->setAddProfilerFunc(reoptimizeIfCallFrequent);
```

Example: do -O2 optimization if function was called more than 10

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static Error reoptimizeTo02(ReOptimizeLayer &Parent, ReOptMaterializationUnitID MUID,
   unsigned CurVerison, ResourceTrackerSP OldRT, ThreadSafeModule &TSM) {
  TSM.withModuleDo([&](llvm::Module &M) {
   auto PassManager = buildPassManager();
   PassManager.run(M);
  });
  return Error::success();
ReOptLayer->setReOptimizeFunc(reoptimizeToO2);
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ReOptLayer->setAddProfilerFunc(reoptimizeIfCallFrequent);

Example: do -O2 optimization if function was called more than 10

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static Error reoptimizeTo02(ReOptimizeLayer &Parent, ReOptMaterializationUnitID MUID,
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```

ReOptLayer->setReOptimizeFunc(reoptimizeToO2); ReOptLayer->setAddProfilerFunc(reoptimizeIfCallFrequent);

DEMO: CLANG-REPL WITH REOPT

- clang-repl is LLVM's in-tree c++ interpreter based on ORC JIT API
- The code originally from CERN's cling which has been used to analyze LHC data.

C JIT API sed to analyze LHC data.

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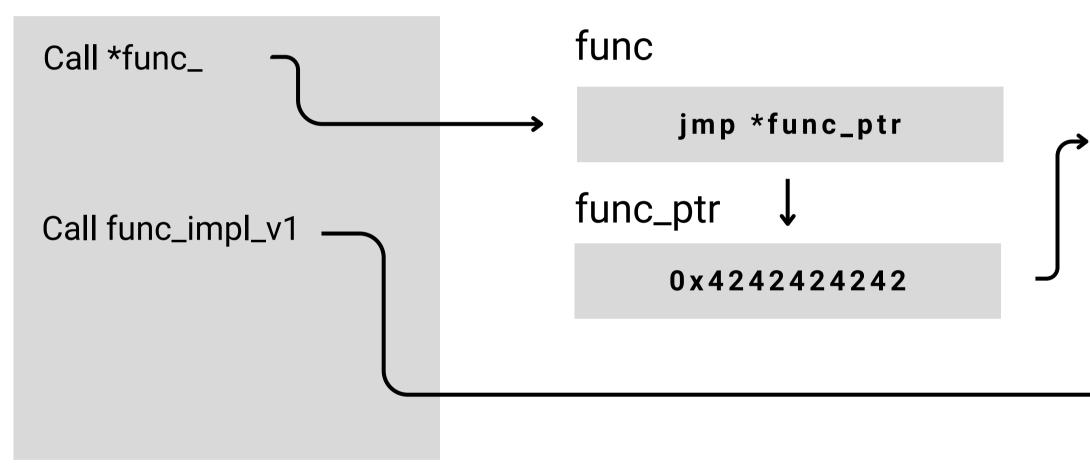
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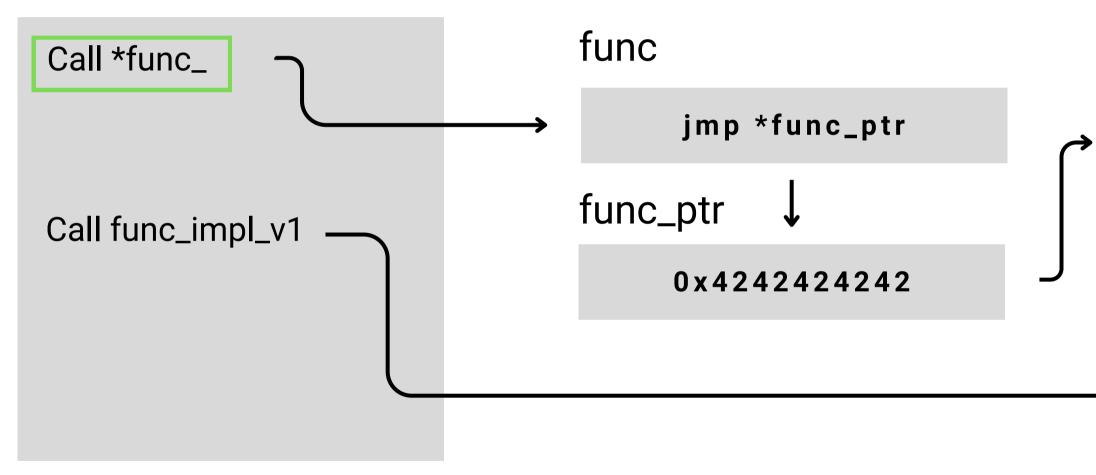
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 - when platform prevents writable and executable memory for security reason.

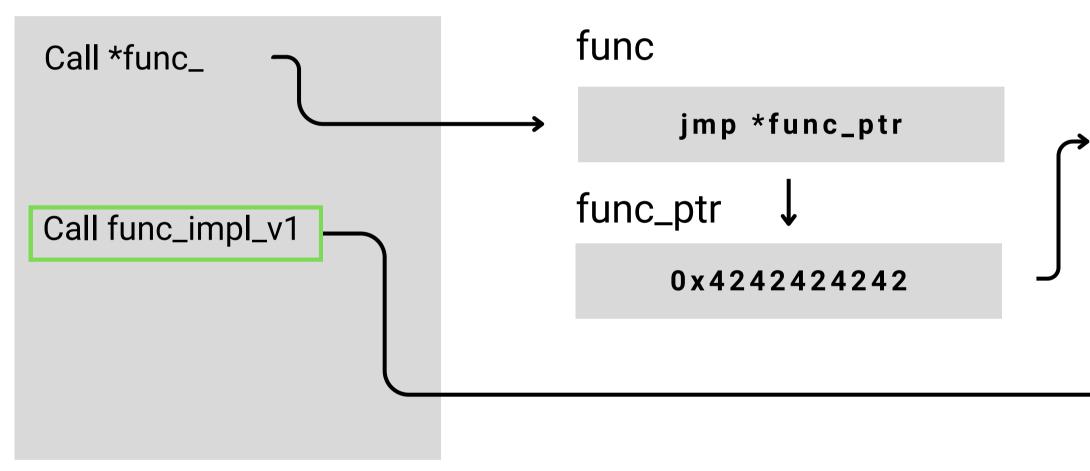
main



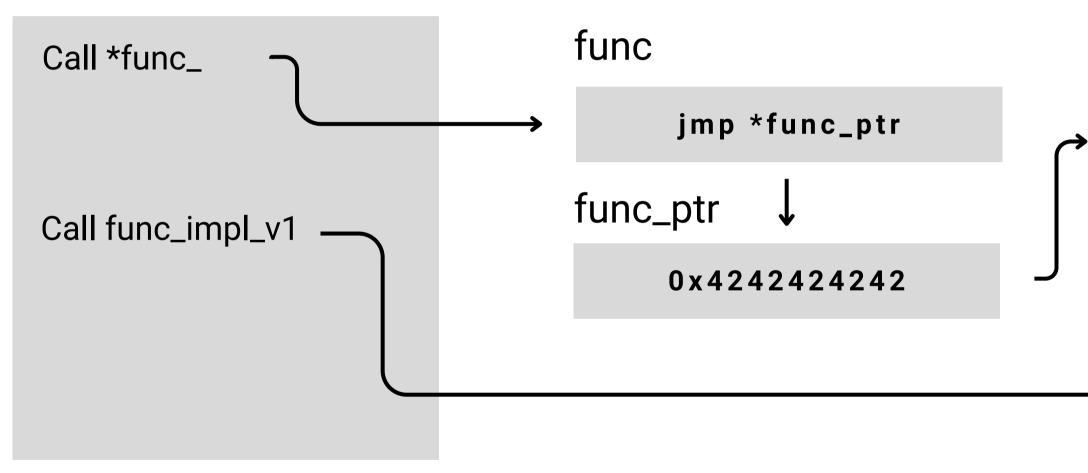
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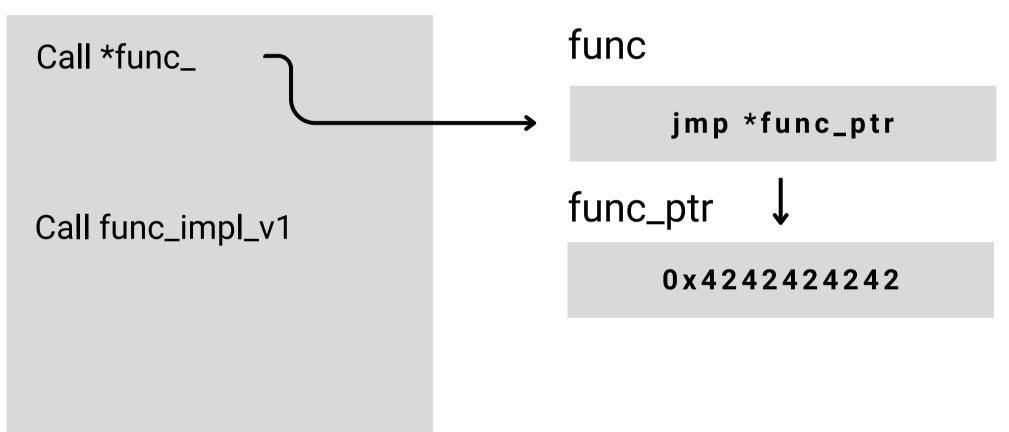
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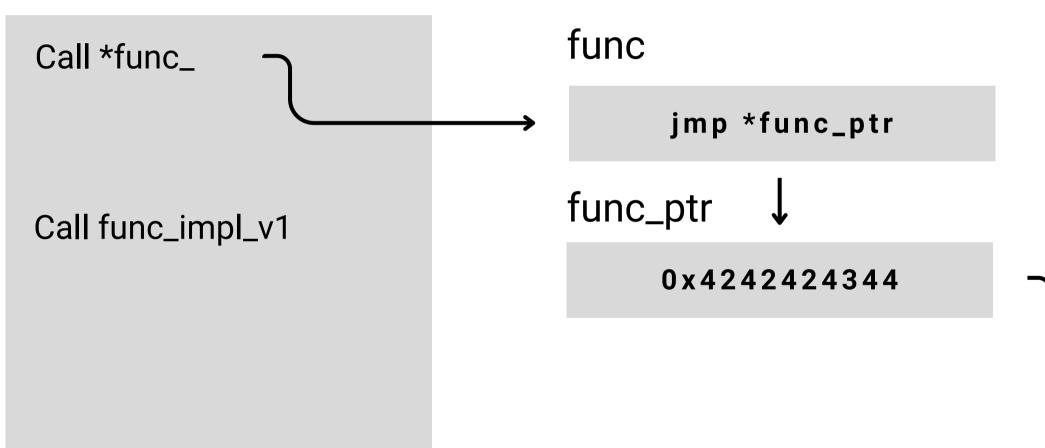
main



func_impl_v1

INTERNALS

main



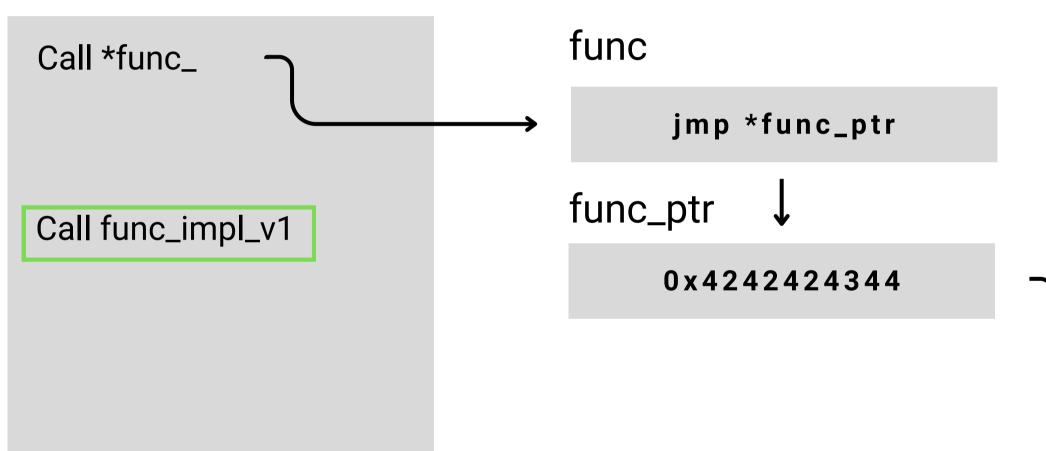
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4

INTERNALS

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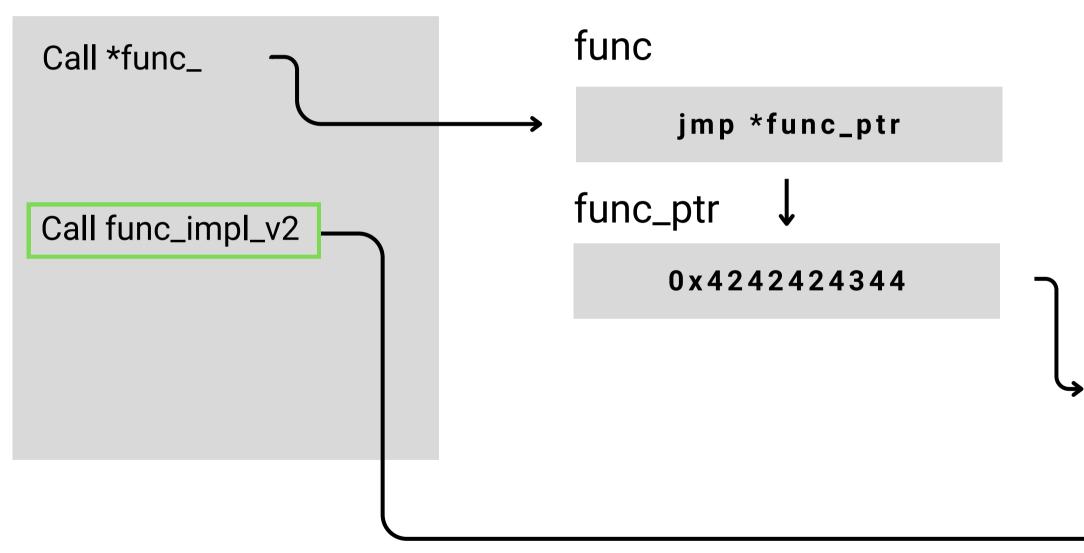
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INTERNALS

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- When calling virtual function, it loads function address from vtable of that classs

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public:
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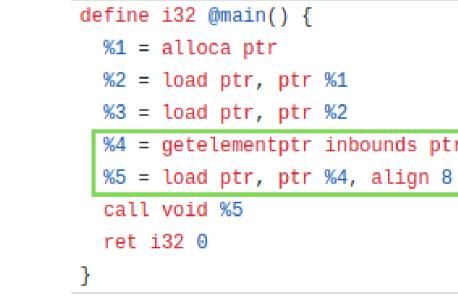


```
define i32 @main() {
  %1 = alloca ptr
  %2 = load ptr, ptr %1
  %3 = load ptr, ptr %2
  call void %5
  ret i32 0
}
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%4 = getelementptr inbounds ptr, ptr %3, i64 0 %5 = load ptr, ptr %4, align 8

Virtual method table

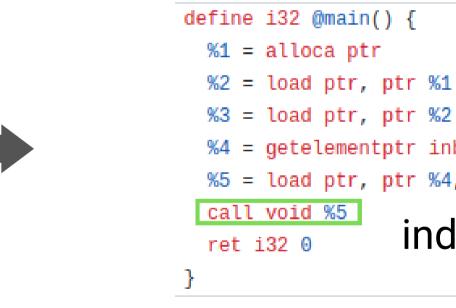
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indirect call

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 - Not just indirection cost but also lose opportunity for potential optimizations as values are not within the same basic block

De-virtualization

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DEMO: CLANG-REPL WITH DEVIRTUALIZATION

• Showcasing the de-virtualization within clang-repl

JIT implementation

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JITted code

call %1 __orc_rt_increment_func_callcnt(%1) __ort_rt_reoptimize(1)

JIT implementation



extern "C"

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JITted code

call %1

orc_rt_reoptimizer.o

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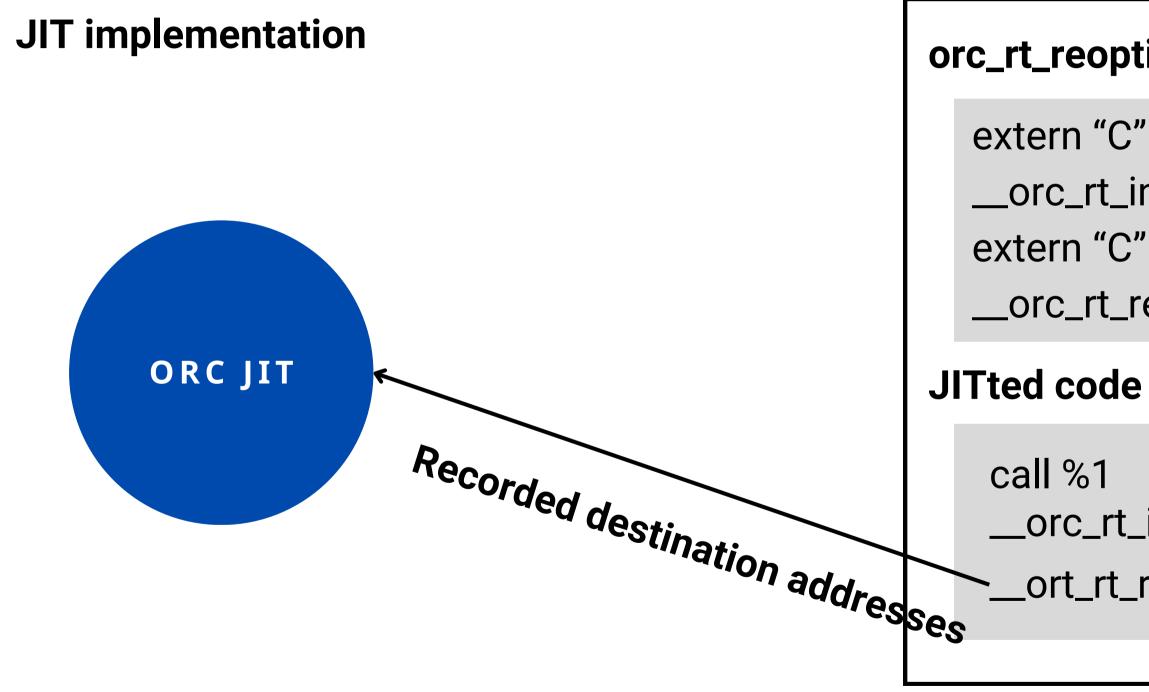
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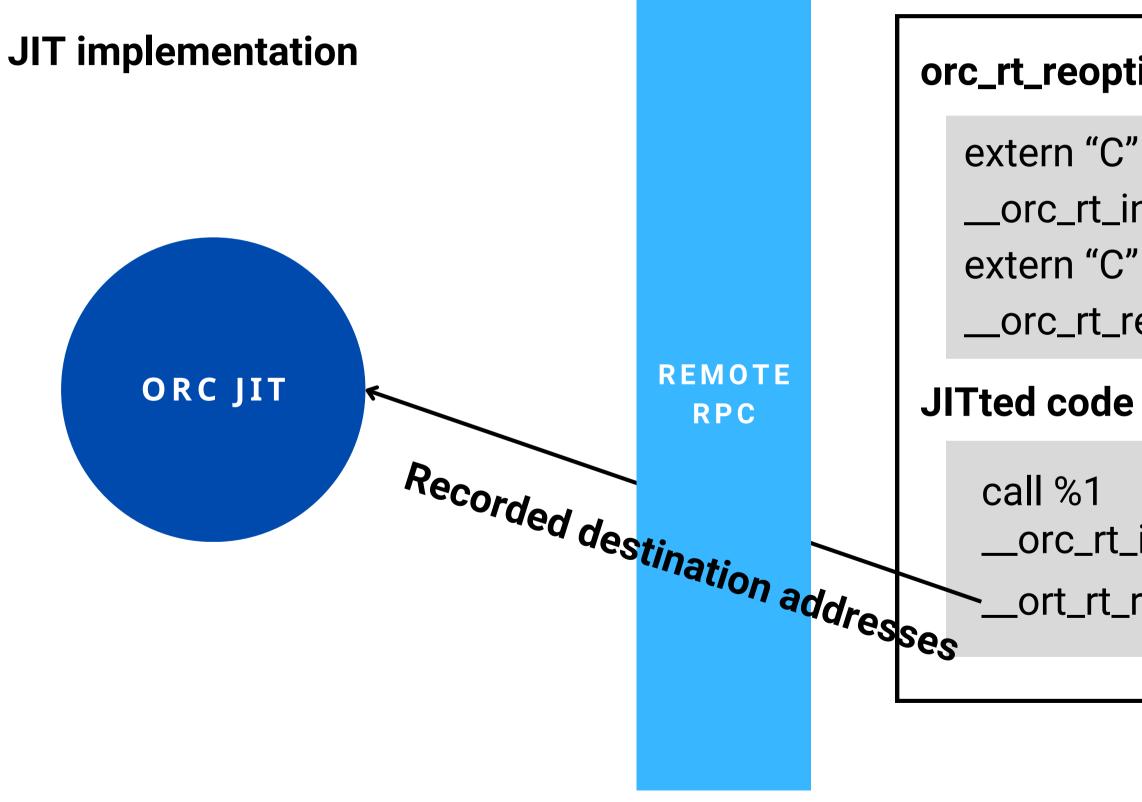


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Program	-01	Reoptimization ON	-02
Boost Spirit (n=1)	1.97s	2.12s	2.24s
Boost Spirit (n=500)	22.46s	21.71s	21.55s

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Program	-00	Reoptimization ON + Devirtualization OFF	-02	Reoptimization ON + Devirtualization ON
Ray Tracer	158.9s	66.6s	66.0s	62.5s

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- The runtime performance drop observed to be as bad as 3x slower.
- Current solution: don't delete function when splitting module but just mark them externally_available.

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- ORC JIT currently have no standard way to inline out-of-module functions.
- Lack of inlining that would have happened in non-reopt mode.
- The runtime performance drop observed to be as bad as 3x slower.
- Current solution: don't delete function when splitting module but just mark them externally_available.
 - $\circ\,$ but this introduces compilation overhead when module is large
 - \circ O(n^2) function duplicates where n is number of functions

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- Penalty = cost for instrumentation + lost optimizations
- Generic JIT profile guided optimization framework
 - Could we possibly overhaul LLVM's existing PGO infrastructure in order to reuse it?

THANKS

Code used today is available at: https://github.com/sunho/LLVM-JIT-REOPT-Example