# Implement and improve an efficient, layered tape with prefetching capabilities

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# The Tape

In clad, the tape is a stack-like data structure that stores intermediate values in reverse mode AD during the forward pass for use during the backward (gradient) pass.

# **Old Implementation**

Clad used a monolithic memory buffer, which, when full while pushing, reallocated double the capacity.

```
template <class InputIt, class NoThrowForwardIt>
CUDA HOST DEVICE void MoveData(InputIt first, InputIt last,
                                NoThrowForwardIt d first) {
  NoThrowForwardIt current = d first;
  for (; first != last; ++first, (void)++current) {
    ::new (const cast<void*>(
    static cast<const volatile void*>(clad addressof(*current))))
      T(std::move(*first));
 constexpr static std::size t init capacity = 32;
CUDA HOST DEVICE void grow() {
  ff (! capacity)
    capacity = init capacity;
     capacity *= 2;
  T* new data = AllocateRawStorage( capacity);
  if (Inew data) {
   destroy(begin(), end());
    printf("Allocation failure during tape resize! Aborting.\n");
    trap(EXIT FAILURE);
```

### Midterm Progress - Slab-based Tape with SBO

- The new tape follows a slab-based structure with small buffer optimization.
- Slab size and buffer size are configurable template parameters.
- Slab size 1024 and buffer size 64 performed the best and were set as default

template <typename T, std::size t 580 SIZE = 64, std::size t SLAB SIZE = 1024> is(T) char raw data[SLAB SIZE \* sameof(T)](); Slab\* next; CUDA HOST DEVICE Slab() : next(nullptr) () CUDA HOST DEVICE T\* elements() { return reinterpret cust<T\*>(raw data); is(T) char m static buffer[SBO SIZE \* SUZEOF(T)][]; bool m using sbo = true; Slab\* m head = nullptr; std::size t m size = 0; CUDA HOST DEVICE T\* sbo elements() { colusplus >= 201703L form std::launder(reluteroret cost<T\*>(m static buffer)): return reinterpret cost<T+>(m static buffer); CUDA HOST DEVICE const T+ sbo elements() const { cplusplus >= 201703L Firm std: launder(reinterpret castsconst T\*>(m static buffer));

values.

```
template <typename... ArgsT>
CUDA HOST DEVICE void emplace back(ArgsT&&... args) {
  if (m size < SBO SIZE) {
   // Store in SBO buffer
::new (const cast<void*>(static cast<const volatile void*>(
        sbo elements() + m size))) T(std::forward<ArgsT>(args)...);
    if (m using sbo)
     m using sbo = false;
     f ((m size - SBO SIZE) % SLAB SIZE == 0) {
      Slab* new slab = new Slab();
      if (Im head) {
        m head = new slab;
        Slab* last = m head;
        while (last->next)
          last = last->next:
        last->next = new slab:
    Slab* slab = m head;
    std::size t idx = (m size - SBO SIZE) / SLAB SIZE:
    while (idx --)
     slab = slab->next;
    ::new (const cast<void*>(static cast<const volatile void*>(
        slab->elements() + ((m size - SBO SIZE) % SLAB SIZE))))
        T(std::forward<ArgsT>(args)...);
 m size++:
```

## Fixes in Benchmarks & Tape Improvements

#### Benchmark Fixes

- Removed Iterations(1) to get better estimate of the benchmarks.
- Fixed Memory Manager Counters.
- Added DoNotOptimize() to prevent compiler from optimizing out the pop function.

#### Tape Improvements

- Added Tail Pointer reducing the push and pop function runs from O(n) to O(1).
- Changed tape to doubly linked list.
- Reused slabs which were previously allocated.

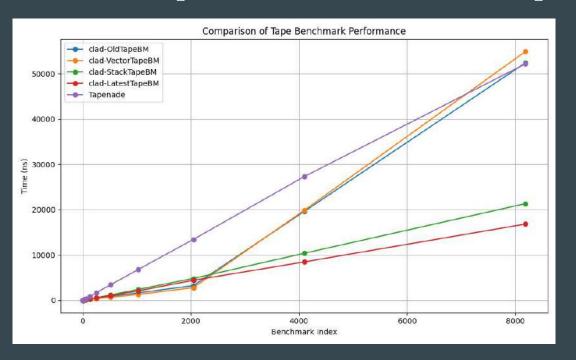
## **Updated Code**

```
template <typename T, std::size t SBO SIZE = 64, std::size t SLAB SIZE = 1024,
         bool is multithread = false>
class tape impl {
 struct Slab {
   alignas(T) char raw data[SLAB SIZE * sizeof(T)];
    Slab* prev;
    Slab* next:
    CUDA HOST DEVICE Slab() : prev(nullptr), next(nullptr) {}
    CUDA HOST DEVICE T* elements() {
#LT cplusplus >= 201703L
      return std::launder(reinterpret cast<T*>(raw data)):
      return reinterpret cast<T+>(raw data);
  alignas(T) char m static buffer[SBO SIZE * sizeof(T)];
 Slab* m head = nullptr;
 Slab* m tail = nullptr;
 std::size t m size = 0;
  std::size t m capacity = SBO SIZE;
```

```
template <typename... ArgsT>
CUDA HOST DEVICE void emplace back(ArgsT&&... args) {
    (m size < SBO SIZE) {
    ::new (const cast<void*>(static cast<const volatile void*>(
        sbo elements() + m size))) T(std::forward<ArgsT>(args)...);
    const auto offset = (m size - SBO SIZE) % SLAB SIZE;
    if (!offset) {
      f (m size == m capacity) {
       Slab* new slab = new Slab():
        (!m head)
         m head = new slab;
         m tail->next = new slab;
         new slab->prev = m tail:
        m capacity += SLAB SIZE;
      ff (m size == SBO SIZE)
        m tail = m head;
       m tail = m tail->next;
    ::new (const cast<void*>(static cast<const valatile void*>(
        m tail->elements() + offset))) T(std::forward<ArgsT>(args)...);
  m size++;
```

#### **New Benchmarks**

• Tested against the old tape, std::vector, std::stack and tapenade.



#### Thread Safety

Added thread-safe tape
 access functions with mutex
 locking mechanism to allow
 concurrent access.

```
CUDACC
template <typename T, std::size t SB0 SIZE = 04, std::size t SLAB SIZE = 1024,
          typename... ArgsT>
T push(tape<T, SB0 SIZE, SLAB SIZE, /*is multithreaded='/true>% to.
       ArgsT&&... val) {
  std::lock quard<std::mutex> lock(to.mutex());
  to.emplace back(std::forward<ArgsT>(val)...);
  return to.back():
template <typename T, typename U, size t N, std::size t SBO SIZE - 64,
          std::size t SLAB SIZE = 1024>
void push(tape<T[N], SBO SIZE, SLAB SIZE, /*is multithreaded=*/true>& to,
               U& val) [
  std::lock guard<std::mutex> lock(to.mutex());
  to.emplace back():
  std::copy(std::begin(val), std::end(val), std::begin(to.back()));
template <typename T, std::size t SB0 SIZE = 64, std::size t SLAB SIZE = 1024>
T pop(tape<T, SBO SIZE, SLAB SIZE, /*is multithreaded="/true>& to) {
 std::lock guard<std::mutex> lock(to.mutex());
 T val = std::move(to.back());
  to.pop back():
 return val;
template <typename T, std::size t N, std::size t SBO SIZE = 64.
          std::size f SLAB SIZE = 1024>
void pop(tape<T[N], SBO SIZE, SLAB SIZE, /*is multithreaded=*/true>& to) {
  std::lock guard<std::mutex> lock(to.mutex());
  to.pop back();
template <typename T, std::size t SB0 SIZE = 64, std::size t SLAB SIZE = 1024>
TG back(tape<T, SBO SIZE, SLAB SIZE, / is multithreaded= /true>& of) {
 std::lock guard<std::mutex> lock(of.mutex());
  return of.back();
```

## Ongoing Progress - Multilayer Tape

#### Multilayer Tape Storage

- Offloading mechanism to offload slabs to disk and load slabs from disk to memory.
- Keep only the last N slabs in memory at a time and offload the rest.
- One slab space where slabs are loaded if element to be loaded is not in memory.

# **Progress**

#### **Current Progress**

- Benchmark script to compare two revisions (PR #1394)
- Modified tape structure to slab-based (PR #1404)
- Added small buffer optimization (PR #1404)
- Enhanced benchmarks (PR #1404 and PR #1533)
- Made improvements to tape(PR #1531 and PR #1633)
- Added thread safety (PR #1470)

## Thank You!